CS 332
Spring 2018

CS 332 - Homework 4

Due: Thursday, March 29

Reading: Chapter 5.1 and 5.3 in the Sipser textbook, pages 213 - 216 and 234-238.
So skip section 5.2 and also those results in 5.1 which refer back to Section 1 of the book.
Chapter 7, pages 275 - 298.

PROBLEMS
Part 1: These 6 problems are to be turned in and each will be graded (10 points each).

1. Recall Dijkstra’s algorithm. This algorithm takes as input an undirected, weighted graph G
with N vertices and a specific vertex s of G. It outputs the length of the shortest path from
s to any other vertex of G.
   If you never learned this algorithm you should have. Look it up in your favorite algorithm
book and see how it works. It’s pretty easy to understand, but go over it and make sure you
see how it works.
   i. What is the complexity (of big-O notation) of Dijkstra’s algorithm? You need not justify
your answer but you should write where you found your version of the algorithm.
   ii. Now consider the “all pairs shortest path” problem. This problem takes the same kind of
input graph G as Dijkstra, but it tell you more. Namely, for each pair of vertices u, v from
G it outputs the length of the shortest path from u to v.
Show how you can reduce the “all pairs shortest path” problem to Dijkstra’s algorithm.
That is, give an algorithm which solves “all pairs shortest path” and which uses Dijkstra’s
algorithm as a procedure call (a subroutine).
   iii. Give a lower bound for the “all pairs shortest path” problem
   That is, find a function f(n) for which it would not be possible to solve the problem in less
than O(f(n)). (I’m not looking for any difficult new algorithm here, just for a “trivial” f which
any algorithm A which solves the problem cannot be smaller than.) Your f should be strictly
smaller than the complexity of the algorithm you found in part ii.

2. Problem 7.8 on page 323 of the Sipser textbook.

3. Say whether each of i. - iv. is True or False. In each case briefly justify your answer.
   i. There are infinitely many undecidable problems.
   ii. If L is decidable and J is undecidable then L ∪ J is undecidable.
   iii. If L is decidable and J is undecidable then L ∩ J is undecidable.
   iv. If J is reducible to K and K is recognizable then J is also recognizable.
4. Prove that the problem of deciding if a TM M ever writes the symbol 1 on its tape is undecidable.
   You should do these problems using reductions.

5. Prove that it is undecidable to determine if two TM’s M and L accept the same languages.
   (Note: So the language which you are to prove undecidable is $S = \{ < M, L > \mid M$ and $L$ accept exactly the same set of strings $\}$.)
   You should do this problem using reductions.

6. You given two languages $J$ and $L$ both of which are sets of natural numbers in P.
   i. Show that $J-L$ is in P
   ii. Show that $JL$ (the concatenation of $J$ and $L$) is in P.

Part 2: These problems are good practice and you should try them. They will not be graded

1. Note: Some of the ideas for this problem were covered in section. You can see the exact definitions you need on page 271 of our textbook.
   a. Give an example of a Boolean formula $F$ which contains at least 2 different variables, and where $F$ is satisfiable and its negation $(\neg)F$ is not satisfiable.
   b. Give an example of a Boolean formula $F$ which contains at least 3 different variables, and where $F$ is satisfiable and has exactly 4 satisfying truth assignments.

2. Page 239, problem 5.7
3. Page 239, problem 5.8